

Secret Message Hiding in WAVE PCM Sound File

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Abstract

Communications today has been done among millions of users using many application devices like internet or satellite communication channels, the services they can transform many varieties of files like text, images, videos, and audio among different places. Therefore the security of data has been of extreme importance in today's information-based society, including the fields of military, diplomacy, corporation, medicine, and etc....

A form of data hiding is steganography, which is contemporary way for protecting the information by embedding data into digital media for the purpose of copyright, and sending secret messages.

In this paper a scheme of steganography system for hiding secret text message in audio file WAV, (Windows Audio Visual) format is proposed, the hiding mechanism was based on using Low-Bit Encoding (LSB) Least Significant Bit substitution techniques.

To support the immunity of the hiding system, encryption methods with some other support methods (i.e., hiding and hopping) were added to the proposed hiding system, a pseudo random number generator has been designed and implemented to generate non-uniform integer jumps between successive hiding events. The jumps mechanism of the suggested generator is based on the linear feedback shift register of length 23 with feedback function $F(x)=1+X+X^{23}$ produce the maximum period $2^{23}-1= 8388607$.

Keywords: ciphertext, cryptography, decryption, encryption , PCM, plaintext, steganography.

إخفاء الرسائل السرية بواسطة الملف الصوتي Wave PCM

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الخلاصة

تتم الاتصالات في الوقت الحاضر من قبل ملايين المستخدمين بواسطة أنواع متعددة من التطبيقات مثل الأنترنت أو قنوات الاتصال عبر الأقمار الصناعية، وهذه الخدمات من خلالها يمكن تحويل العديد من أنواع الملفات مثل النصوص، الصور، وأشرطة الفيديو والرسائل الصوتية بين مواقع مختلفة من الكرة الأرضية. وبالتالي فإن أمن البيانات وسريتها لها أهمية بالغة وتتطلب عناية قصوى أثناء تداولها في مجتمعات اليوم التي تعتمد على المعلومات والبيانات أساساً لها في المجالات العسكرية، الدبلوماسية، الشركات التجارية والمؤسسات الطبية وحتى على مستوى الأفراد.

اقترح نظام إخفاء النصوص السرية في صيغة الملفات الصوتية WAV (Windows Audio Visual) وقد اعتمدت آلية الإخفاء بطريقة (LBE) Low-Bit encoding أو تقنيات الاستبدال (LSB) Least Significant Bit.

ولدعم حصانة نظام الإخفاء هذا تم إضافة طرق تشفير أخرى encryption مع أساليب الدعم الأخرى مثل (الأختباء والتنقل)، أضيفت هذه الطرق لدعم نظام الإخفاء المقترح وتم تصميم مولد رقم عشوائي زائف لتوليد عدد صحيح غير موحد بين أحداث الإخفاء المتعاقبة.

إن ميكانيكية القفز لمولد عشوائي مقترح يعتمد على مسجل التزحيف ذو التغذية الخلفية الخطية (LFSR) بطول 23 وبدالة تغذية خلفية $F(x)=1+x+x^{23}$ بحيث تولد أعلى قيمة $2^{23}-1=8388607$

1- Introduction

Steganograph (literally, covered writing) is the hiding of secret messages within another seemingly innocuous message, or carrier (cover). Digital carriers include email, audio, and video messages, disk space, disk partitions, and images [1].

The subjective quality of the audio data depending on our hearing sense could not recognize all voices and noises that are accompanied with original wave media. Data hiding in audio signals are especial challenge, because the (HAS) Human Auditory System operates over a wide dynamic range. The HAS perceives over a range of power greater than billion to one and range of frequencies greater than one thousand to one.[2]

Sensitivity to additive random noise is also acute. When performing data hiding on audio, one must exploit the weakness of the HAS, while at the same time being aware of extreme sensitivity of the human auditory system. [3]

In this research has been the builder a hiding system of secret messages in wave files without producing any significant distortion

2 –Least Significant Bit Encoding (LSB)

Low-bit encoding is the simplest way to embed data into other data structure. By replacing the least significant bit of each sampling point by a code binary string, we can encode a large amount of data in an audio signal. The bit rate will be 8 Kbps in an 8KHz sampled sequence and 44Kbps in a 44KHz sampled sequence.[4]

The simple algorithm of LSB for low-bit encoding is as follows :

- Represent the object as vector of integers.
- Change the least significant bit in either all or some integers to represent a 1 or 0 in the mark

Depending on the amount of embedded information and the amount when the wave media carries, it is quite unperceivable. For example the first thing to do is to hide an ASCII code of A, which represent 01000001, and then part of wave media cover data will be

(11101011 10001011 10101010 11001010 10001010 10111101 11111111 01001111)

After information is hidden the wave file will be

(11101010 10001011 10101010 11001010 10001010 10111100 11111110 01001111)

3 - Wave PCM Sound file Format

The most popular waveform coding technique used to present the human speech using Pulse Code Modulation (PCM).[3]

The WAVE file format is a subset of Microsoft’s RIFF specification for the storage of multimedia files. A (RIFF) resource enter change file starts out with a file header followed by a sequence of data chunks as shown in figure (1). A WAVE file is often just a RIFF file with a single “WAVE” chunk which consists of two sub-chunks—a “fmt” chunk specifying the data format and a “data” chunk containing the actual sample data which is the “canonical form”. As shown on table. (1)

File offset (byte)	Field name	Field size (byte)	The RIFF chunk descriptor
0	Chunk ID	4	The “RIFF” chunk descriptor The format of concern here is “WAVE”, which requires two sub-chunks: “fmt” and “data”
4	Chunk size	4	
8	Format	4	
12	SubChunk1ID	4	The “fmt” sub-chunk Describes the format of the sound information in the data sub-chunk
16	SubChunk1size	4	
20	Audio Format	4	
22	Num Channels	2	
24	Sample Rate	2	
28	Byte Rate	4	
32	Block Align	2	
34	Bit spersample	2	
36	SubChunk2ID	4	The “data’ sub-chunk Indicates the size of the sound information and contains the raw sound data
40	SubChunk2size	4	
44	Data	Subchunk2size	

Figure (1) The Canonical WAVE file format [4]

4 - Linear Feedback Shift Register

A commonly used method for generating binary sequences, especially pseudo-random sequences, is to feed a binary function of the state of a shift register back to its input. The stages of shift register serve as the input of logical circuit whose output is connected to the input of the shift register. Shown in figure (2)

An n-stage shift register s_0, s_1, \dots, s_{n-1} . The contents of the stages change in time with a clock pulse according to the rule:

Let $S_i(t)$ denote the content of s_i after the t^{th} time pulse ($t=0,1,2,\dots$).

$S_{n-1}(t+1)=f(c_0s_0(t), c_1s_1(t),\dots c_{n-1}s_{n-1}(t))$, where the c_i are all specified as 0 or 1.

The function build is called the feedback function of the register and if $f(s_0, s_1, \dots, s_{n-1}) = c_0s_0 + c_1s_1 + \dots + c_{n-1}s_{n-1}$ then the register called linear register. This is represented by the $c_i=1$ denotes a closed connection and $c_i=0$ an open one.[6]

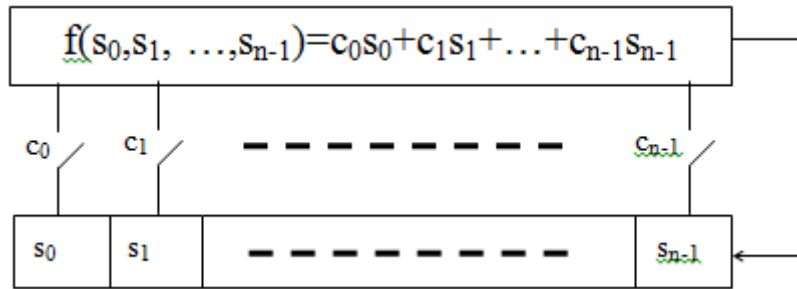


Figure (2) : Linear feed back shift register

A assume $c_0=1$ so that $s_{n-1}(t+1)$ is dependent on $s_0(t)$. Let $s_t=s_0(t)$, an infinite binary sequence denoted s_t satisfy the linear recurrence relation as $s_{t+n} = \sum_{i=0}^{n-1} c_i s_{t-i}$, for $t=0,1,\dots,n$.

The shift register has been identified as characteristic polynomial as shown [8].

$$F(x) = 1 + c_1x + c_2x^2 + \dots + c_{n-1}x^{n-1} + x^n \text{ (remembering } c_0=1\text{)}.$$

5 - Cryptography

Steganography is a way that deals with finding the best place in cover media to hide data. If the data encrypted before hiding it, this will give more security immunity to the cover data.

In this paper research, before hiding the secret data were encrypted by using stream cipher generator. [7]

5.1 - Stream Cipher

One of the cryptographic primitives used to ensure secure communication over public and unsecured channels (such internet, mobile) is the stream cipher. In a stream cipher the plaintext is encrypted on bit by bit basis. In the encrypting of data flow transmitted, the key is fed into an algorithm called running key generator (RKG) to generate a long pseudorandom binary sequence. This "Key Stream" is then mixed with the plaintext sequence, usually by using exclusive-or (XOR bitwise module 2 additions) logic gate, to produce the cipher text. A typical stream cipher is shown in figure (3)

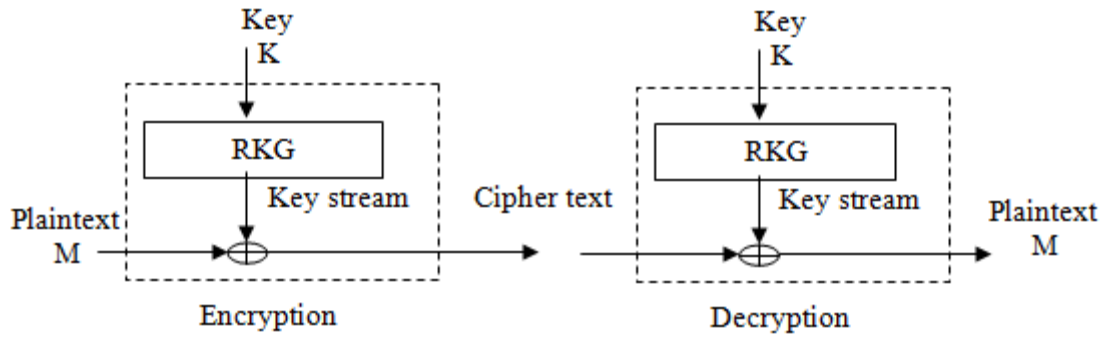


Figure (3) Stream Cipher

A common type of RKG employed in stream cipher system consists of n (mostly maximum length) LFSRs whose output sequences are combined in a nonlinear function F to produce the key stream.[7]

5.2- Encryption and Decryption

Cryptographic algorithm have been carefully designed for maximum security, it divide into two subsystem:

Driving subsystem, which are consists of 8-LFSRs with maximum period are corresponding feedback polynomials

$$\begin{aligned}
 F_1(X) &= 1+X^{13}+X^{33} & F_2(X) &= 1+X^3+X^{31} \\
 F_3(X) &= 1+X^2+X^{29} & F_4(X) &= 1+X^3+X^{28} \\
 F_5(x) &= 1+X^3+X^{25} & F_6(X) &= 1+X^5+X^{23} \\
 F_7(x) &= 1+X^{14}+X^{17} & F_8(X) &= 1+X^{14}+X^{17}+X^{18}+X^{19}
 \end{aligned}$$

The second subsystem is non-linear compost combining subsystem F. which represent a 2D matrix 16x8 bits each bit addressed by driving subsystem, the output determined through the intersection of row by first 4 LFSRs (0..15 row) and column by next 3 LFSRs (0..7 row) the result mixed with the output of the register 8 using x or to produce the key sequence, as shown in figure (4)

6 -

original code 6bit			b_5	b_4	b_3	b_2	b_1	b_0
	b_5	b_4	b_3	b_2	b_1	A_0		
	b_5	b_4	b_3+b_5	b_2+b_4	b_1+b_3	$b_0 + b_2$	b_1	b_0
New extended code 8bit	a_7	a_6	a_5	a_4	a_3	a_2	a_1	a_0

Cryptographic Algorithm

Build a coding table of 6-bit has been build of printable keyboard characters as in the coding table(1). The basic key is 16-character of 6-bit extended into 8-bit by using the original character shift left 2 bits and x.

Example:

After extended the basic key into 128 bits we use this bits to initial the driving and the nonlinear combining parts, each shift of driving subsystem will produce one bit output from combining and then to produce key bit.

original code 6bit			b_5	b_4	b_3	b_2	b_1	b_0
	b_5	b_4	b_3	b_2	b_1	A_0		
	b_5	b_4	b_3+b_5	b_2+b_4	b_1+b_3	$b_0 + b_2$	b_1	b_0
New extended code 8bit	a_7	a_6	a_5	a_4	a_3	a_2	a_1	a_0

dec	code	binary	dec	code	Binary	Dec	code	binary	dec	code	binary
0	S	000000	16	T	010000	32	f	100000	48	3	110000
1	q	000001	17	9	010001	33	C	100001	49	d	110001
2	t	000010	18	/	010010	34	8	100010	50	J	110010
3	p	000011	19	m	010011	35	7	100011	51	H	110011
4	x	000100	20	Q	010100	36	O	100100	52	k	110100
5	g	000101	21	c	010101	37	n	100101	53	L	110101
6	0	000110	22	y	010110	38	G	100110	54	K	110110
7	w	000111	23	u	010111	39	X	100111	55	E	110111
8	R	001000	24	i	011000	40	Z	101000	56	6	111000
9	o	001001	25	b	011001	41	5	101001	57	V	111001
10	B	001010	26	M	011010	42	D	101010	58	e	111010
11	l	001011	27	v	011011	43	P	101011	59	r	111011
12	z	001100	28	w	011100	44	I	101100	60	space	111100
13	A	001101	29	F	011101	45	a	101101	61	h	111101
14	4	001110	30	1	011110	46	s	101110	62	j	111110
15	Y	001111	31	N	011111	47	2	101111	63	U	111111

Table (1) Cryptographic Algorithm

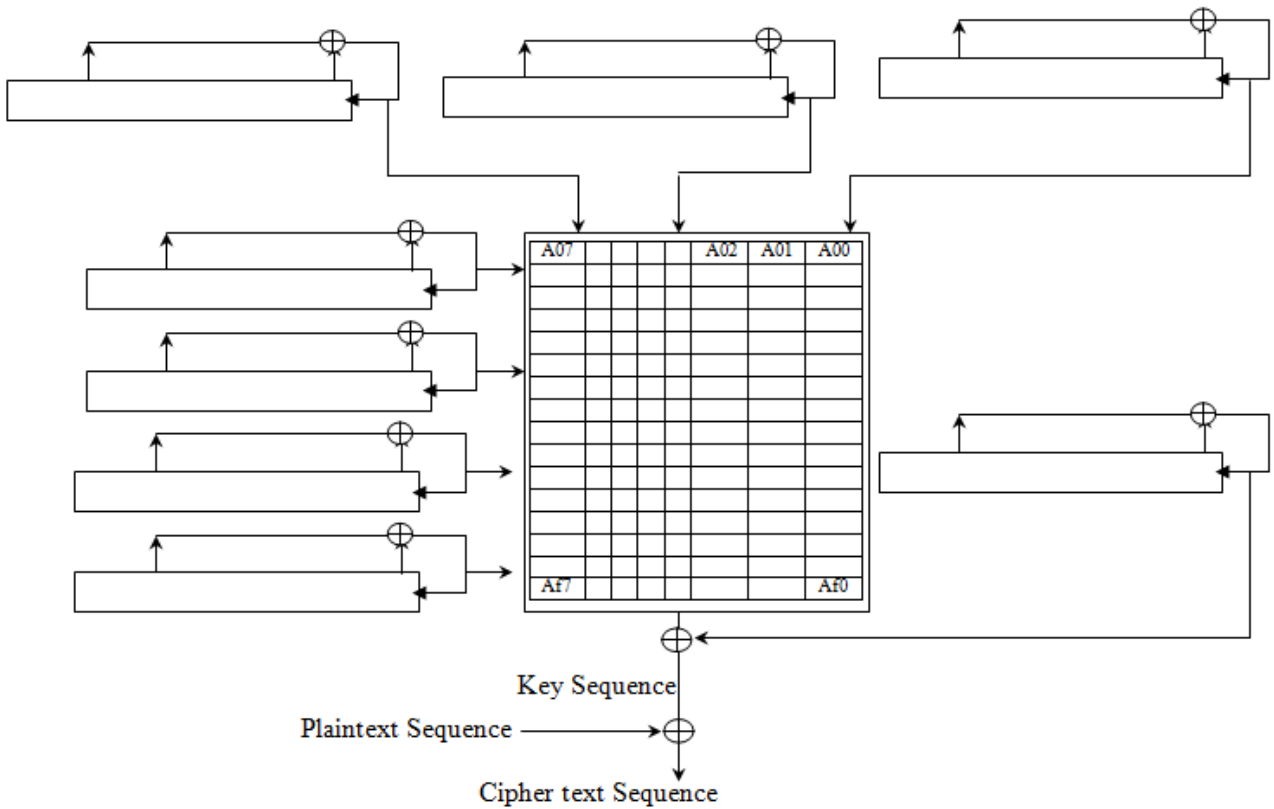


Figure (4) Cryptographic mechanism

7 - Hopping and Hiding

If the data embedding is done sequentially over the entire wave file (i.e. hiding in byte after byte), will make the stego system very vulnerable against simple tools of steganalysis. Therefore jump with variable (pseudo-random) length between successive hiding events will greatly increase the security level of stego system. In this research a pseudo random number generator was designed and implemented to generate non-uniform integer jumps between successive hiding events. The jumps mechanism of the suggested generator is based on the linear feedback shift register of length 23 with feedback function $F(x)=1+X+X^{23}$ produce the maximum period $2^{23}-1=8388607$. [9]

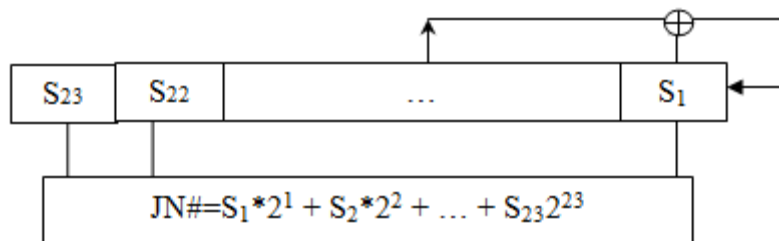


Figure (5) The jumps mechanism

The Algorithms of the proposed system :-

Hiding poses

Step 1 : read plan text

Step 2: split play text into an ASCII

Step 3 : covered each ASCII info eight bits

Step 4 : read the decimal six, bits key

Step 5 : shift key char 2 bit

Step 6 : hopping the eight bits with R key G inside audio cover

Step 7 : hide the cipher text bit in the LSB of determined Audio byte

Step 8: if plan file finished

 go to step 9 etc

 go to step 6

extracting process

step 9 : hopping inside Audio file (cover)

step 10 : mask with LSB of determined audio byte and extract the cipher text bit

step 11 : extract hex decimal key of six bits code

step 12 : shift key char 2 bit left then mixed it with the original by using XOR

step 13: collect bit into ASCII code

step 14 : if cipher text finished

 go to step 15 etc

 go to step 9

step 15 : end

8 - The proposed system

The block diagram explains how to hide the message and how to extract it shown in figure (6)

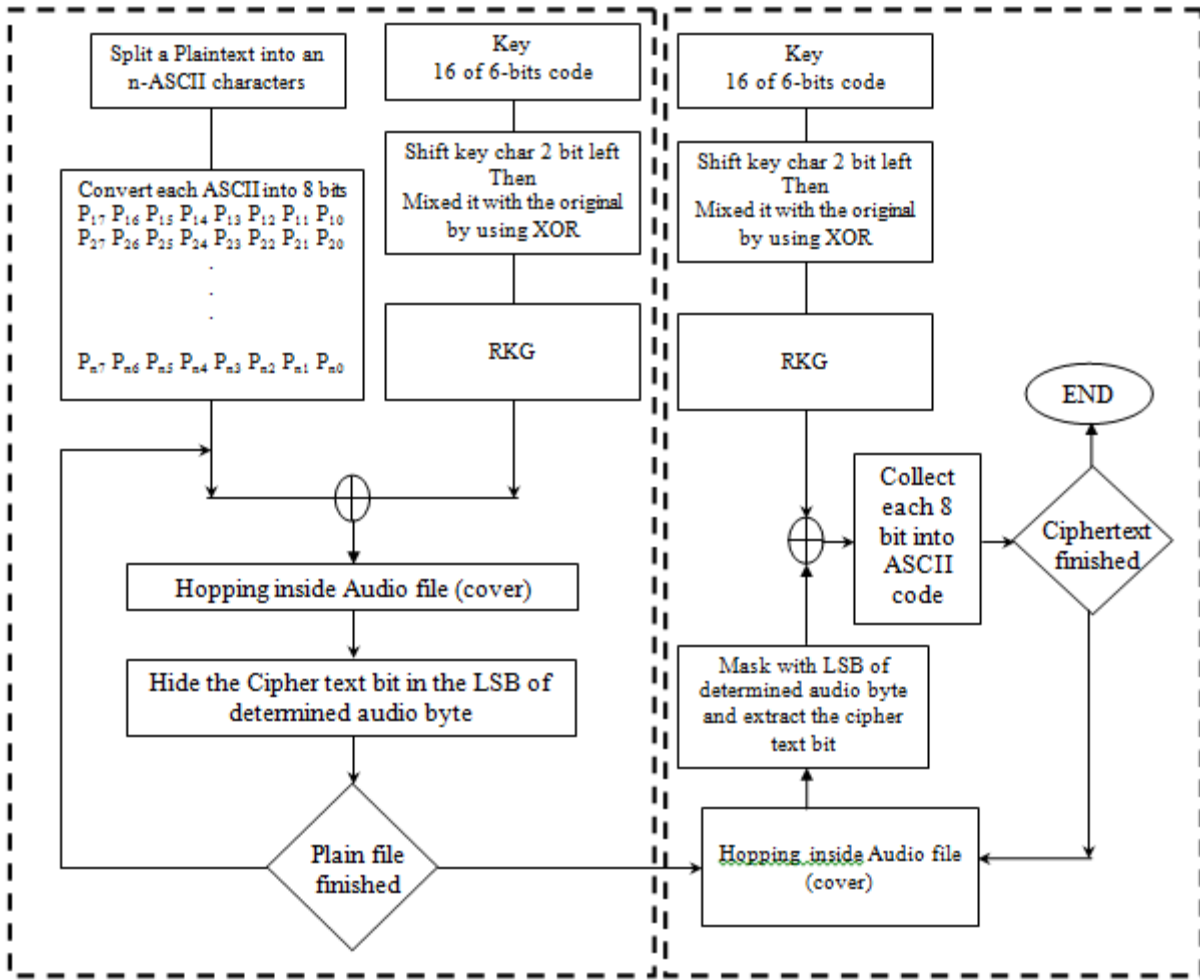


Figure (6) The block diagram of proposed system

9 - Fidelity Criteria

Signal-to-noise (SNR) measures are estimates of the quality of a quality reconstructed image compared with an original image. Reconstructed images with higher metrics are judged better quality. Traditional SNR measures do not equate.

First compute the mean squared error (MSE) of the reconstructed audio as follows:

$$MSE = \text{MSE} = \frac{\sum [f(I,j) - f^*(I,j)]^2}{N^2} \dots\dots\dots(1)$$

The summation is over all signals. The root mean square error (RMSE) is the square root of MSE.

PSNR in decibels (dB) is computed by using:

$$PSNR = 20 \log_{10} \frac{(L-1)^2}{RMSE} \dots\dots\dots(2)$$

L: Max Value

$$20SNR = 20 \log_{10} \left(\frac{\text{Max value}}{RMSE} \right) \dots\dots\dots(3)$$

Mean absolute error (MAE) can measure the quality to the different of a reconstructed audio compared with one original audio, the value of this measure be between 1 and 0, the actual value be good if the value near from zero.

$$MAE = \text{Bit Per Sample BPS} = \frac{\text{Hidden data Bit}}{\text{Total Cover size (sample)}} \dots\dots\dots(4)$$

Typical PSNR values range will be greater than or equal to 50 .They are usually reported to two decimal points (e.g 25.47). The actual value is not meaningful, but the comparison between two values for different constructed audio signals gives one measure of quality. An informal threshold of 0.5 dB PSNR is used to decide whether to incorporate a coding optimization because they believed that an improvement of that magnitude would be visible.

The data hiding in wave data, and samples, the following table (2) illustrates the PSNR, MSE, and BPS results for eight types of data with different sizes are hidden in a “boop.wav” file whose size is (79561).

“boop” file whose size is (64162 byte).					
File	Length(byte)	MSE	PSNR	BPS	SNR
1	1925	0.01	68	0.19	66
2	3208	0.02	65	0.32	64
3	6416	0.04	62	0.64	61
4	8341	0.05	61	0.83	59
5	10266	0.07	60	1.03	58
6	12832	0.08	59	1.29	57
7	17324	0.1	58	1.74	55
8	19249	0.11	58	1.93	55

Table (2) : Output result of MSE,BPS, SNR and PSNR

The general structure of the proposed system is illustrated in figure (6) it consists of two basic modules: hiding and extraction modules . The input to this system are the cover file (wave file) and secret file (binary file) . These input are processed in the hiding part with various operations to produced stego wave file . The stego audio entered to extraction stage is processed through a set of operations to retrieve the secret data.

10- Conclusions:-

From the test results listed in propose system the following remarks wave derived

- 1- Hiding in voiced block sample is more suitable to avoid noise occurrence which is more probably happen when unvoiced blocks are used as host area.
- 2- Large threshold value provide more power in cover audio signal by avoiding unvoiced blocks and increased correct retrieved bits but decreased in hiding.
- 3- The results show acceptable hiding performance and the quality of reconstructed wave file is not subjectively different from the original wave.

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